

CSE560 Midterm Exam

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Syntax of Datalog (8 Marks)

Question 1 (2 marks)

Consider the constants *tom* and *sally*, and the 0-ary predicate *raining* and the 2-ary predicate *mother*(-, -). What are all of the legal **atomic** clauses (also called facts) that can be written.

Answer:

```
raining
mother(tom, sally)
mother(sally, tom)
mother(tom, tom)
mother(sally, sally)
```

Question 2 (2 marks)

Give two non-atomic clauses (also called rules) that are legal.

Answer:

Many possible answers. Here are two:

```
raining <- mother(tom, sally)
mother(tom, sally) <- mother(sally, tom)
```

Question 3 (2 marks)

If we add the function *motherof*(-), give an *atomic* clause that is now legal.

Answer:

Many possible answers. Here is one:

```
mother(motherof(tom), tom).
```

Question 4 (2 marks)

With the addition of the function *motherof*(-), how many different atomic clauses can we now write?

Answer:

infinite number, as we can have arbitrary embedding of *motherof*.

Intended Interpretation (4 Marks)

Question 5 (2 marks)

Annie is Beth's mother, and Beth is Corine's mother. Imagine that you are a knowledge engineer and have the intended interpretation where a maps to Annie, b maps to Beth, and c maps to Corine, and $mother(X,Y)$ is true if X is Y 's mother.

The following two clauses are clearly true of the intended interpretation.

$mother(a,b)$

$mother(b,c)$

Mark each of the following clauses as to whether it is true or false in the intended interpretation.

$mother(c,a)$

Answer:

false

$mother(b,a)$

Answer:

false

$mother(a,b) \leftarrow mother(b,c)$

Answer:

true, since both parts are true.

Note that ' \leftarrow ' does not capture any sense of causality.

$mother(a,b) \leftarrow mother(c,a)$

Answer:

true, as true \leftarrow false is always true.

$mother(c,a) \leftarrow mother(a,b)$

Answer:

false, as false \leftarrow true is always false.

Question 6 (2 marks)

Now say we add the function $motherof(X)$ that will indicate who the mother of X is.

Write down one fact that the knowledge engineer should add to the knowledge base to capture the full meaning of $motherof$.

Answer:

$mother(motherof(X), X)$

Interpretations (11 Marks)

Question 7 (2 marks)

Consider the language with constants a and b , and the 0-ary predicate $raining$ and the 1-ary predicate $girl(_)$. For the domain $D = \{x, y, z\}$, how many different ways can an interpretation map the constants (how many different ϕ 's are there)? Explain your answer.

Answer:

$a \rightarrow x, y$ or z

$b \rightarrow x, y$ or z

So, 3×3

Question 8 (1 marks)

How many ways can an interpretation map the predicate *raining* (how many different ways can π map *raining*)?

Answer:

raining \rightarrow T/F 2 ways

Question 9 (2 marks)

How many ways can an interpretation map the predicate *girl*($_$). Give one of the mappings in full.

Answer:

girl(x) \rightarrow T/F

girl(y) \rightarrow T/F

girl(z) \rightarrow T/F

so $2 \cdot 2 \cdot 2$

girl(x) \rightarrow T, girl(y) \rightarrow T, girl(z) \rightarrow T

Question 10 (1 marks)

How many interpretations are there altogether?

Answer:

$3 \cdot 3 \cdot 2 \cdot 2 \cdot 2$

Question 11 (1 marks)

If we add the variable *i* to our syntax, does this change the number of interpretations? Why or why not?

Answer:

No. Interpretations do not specify the mapping of variables, that is what rho does.

Question 12 (1 marks)

How many of the interpretations from Question 8 are models of $KB = \{\}$.

Answer:

All of them

Question 13 (3 marks)

How many of the interpretations are models of $KB = \{girl(a), girl(b), raining\}$

Answer:

Case 1: $\phi(a) = \phi(b)$

3 different ϕ 's

Without loss of generality, say $\phi(a) = x$

then girl(x) must be true,

and does not matter whether girl(y) and girl(z) is true or false

raining must map to true

So, $2 \cdot 2$ ϕ 's

Case 2: $\phi(a) \neq \phi(b)$

6 different ϕ 's (3 ways to map a and 2 ways to map b)

Without loss of generality, say $\phi(a) = x$ and $\phi(b) = y$

then girl(x) and girl(y) must be true.

and does not matter whether girl(z) is true or false

raining must map to true

So 2 ϕ 's

Altogether $3 \cdot 2 \cdot 2 + 6 \cdot 2$ models

Semantic Proofs (9 Marks)

Question 14 (2 marks)

What does $KB \models \alpha$ mean? Phrase your answer in terms of interpretations (not models).

Answer:

For any interpretation I that makes KB true, I also makes α true.

Question 15 (1 marks)

What does it mean if α does not logically follow from KB ? Does this mean that α is false in our intended interpretation? Explain your answer.

Answer:

It means there is at least one model of KB that makes α false. But, there might be models of KB that make α true. So, the intended interpretation might make α true or false.

Question 16 (3 marks)

Consider the following KB

p .

$q \leftarrow p$.

Use a truth table to prove that q logically follows from KB ($KB \models q$). Make sure you explain why your truth table shows $KB \models q$.

Answer:

p	q	$q \leftarrow p$	Model
T	T	T	yes
T	F	F	no
F	T	T	no
F	F	T	no

Only model is p and q are both mapped to T. So, in all models, q is true. So, q follows from KB .

Question 17 (3 marks)

Say that a is a constant in our syntax and our KB consists of the following.

$p(X)$.

Give a semantic proof that $p(a)$ logically follows from KB . (Hint: do a proof by contradiction involving interpretations, models, and/or variable assignments.)

Answer:

Assume that $p(a)$ does not follow from KB . So, there must be an interpretation I that is a model of KB but $I(p(a))$ is false.

Let I consist of ϕ, π and D .

Say $\phi(a)$ maps to the domain object foo . So, $\pi(p)(foo)$ must be false.

Let ρ be a variable assignment that maps X to foo . Since I is a model of KB , and $p(X)$ in KB , thus, I and any variable assignment, including ρ , must make $p(X)$ true.

So, ρ and σ must make $p(X)$ true.
 So, $\pi(p)(foo)$ must be true.
 Contraction.

So $p(a)$ must follow from KB

Derivations (8 Marks)

Question 18 (2 marks)

What is the most general unifier of the following.

$p(a,b)$ $p(Y,Z)$

Answer:

Y/a Z/b

$p(a,b,X)$ $p(Y,Z,Z)$

Answer:

Y/a Z/b X/b

$p(q(X),q(a))$ $p(Z,Z)$

Answer:

$Z/q(a)$ X/a

Question 19 (3 marks)

Consider the following KB.

- 1 $parent(X, Y) \leftarrow mother(X, Y)$
- 2 $parent(X, Y) \leftarrow father(X, Y)$
- 3 $ancestor(X, Y) \leftarrow parent(X, Y)$
- 4 $ancestor(X, Z) \leftarrow parent(X, Y) \wedge ancestor(Y, Z)$
- 5 $mother(amy, bob)$

Give a top-down derivation (from query to KB) of $ancestor(amy, bob)$. Show the details of the derivation (which rule used, and the substitution).

Answer:

```
yes <- ancestor(amy, bob)
  use: ancestor(X, Y) <- parent(X, Y)
  sub X/amy Y/bob
yes <- parent(amy, bob)
  use: parent(X, Y) <- mother(X, Y)
  sub X/amy Y/bob
yes <- mother(amy, bob)
  use mother(amy, bob)
yes <-
```

Question 20 (2 marks)

Give a bottom-up derivation. Show all steps.

Answer:

$mother(amy, bob)$	KB	1
$parent(amy, bob)$	1 and $parent(X, Y) \leftarrow mother(X, Y)$ with X/amy Y/bob	2

ancestor(amy,bob) 2 and ancestor(X,Y) <- parent(X,Y) with X/amy Y/bob 3

No other facts can be added.
ancestor(amy,bob) is true since in C.

Question 21 (1 marks)

Since the computer only does derivations, why is it necessary to have semantics?

Answer:

We need to make sure that the derivations that the computer makes are true and that it is finding all possible derivations. We use semantics to check this.

Defining Knowledge (4 Marks)

Question 22 (2 marks)

Using definite clauses (rules and facts), define **pulloutone(Item,List,Remainder)**, which is true if **Item** is in **List**, and **Remainder** is the **List** minus **Item**. You can use either the prolog list notation $[H|T]$ (with $[]$ as the empty list) or the function notation $\text{cons}(H,T)$ (with nil as the empty list).

Answer:

```
pulloutone(I, [T|Rest], Rest).  
pulloutone(I, [T|Rest], New) <-  
    putoutone(I, Rest, Remainder)  
    New = [T|Remainder]
```

Question 23 (1 marks)

What will the answer be for the following query? (You do not have to give a derivation; just give the answer.)

```
?pulloutone(p(X), [p(a), p(b), r, q(c), Z, p(d)], L).
```

Answer:

```
X=a L=[p(b), r, q(c), Z, p(d)]
```

Question 24 (1 marks)

What are the other answers (if we did an exhaustive search for all of the possible answers)?

```
?pulloutone(p(X), [p(a), p(b), r, q(c), Z, p(d)], L).
```

Answer:

```
X=b      L=[p(a), r, q(c), Z, p(d)]  
Z=p(X)   L=[p(a), p(b), r, q(c), p(d)]  
X=d      L=[p(a), p(b), r, q(c), Z]
```

Search (2 Marks)

Question 25 (1 marks)

What is the role of search in converting a top-down proof procedure into a reasoning procedure.

Answer:

Search is how the non-determinism of the top-down proof procedure is resolved. It ensures that all choices for what rule are used to resolve with the first atom of the answer clause is eventually searched.

Question 26 (1 marks)

Does search play a role in converting a bottom-up proof procedure into a reasoning procedure? Why or why not.

Answer:

No. The bottom-up proof procedure works by finding all consequents from a knowledge base. So, it already does search for all possibilities.
